# OPER 627: Nonlinear Optimization Lecture 16: Algorithms for optimization over a simple set 

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## Today's Outline

(1) Constrained gradient descent: conditional gradient method (Frank-Wolfe)
(2) Gradient projection method: an iterative algorithm for solving constrained optimization over a simple set
(3) Proximal point: extended gradient projection

## Quiz

(1) What is the optimality conditions for constrained optimization problems over a closed convex set
(2) What is the optimality conditions for constrained optimization problems over the nonnegative orthant $\mathbb{R}_{+}^{n}$ ?
(3) What is a projection? What are the properties of a projection operator?

## Constrained gradient descent algorithms

Algorithms in this lecture:
(1) They do not rely on any structure of the constraint set other than convexity
(2) They generate sequences of feasible points by search along descent directions

Algorithm ingredient:

- Feasible direction $d \neq 0: x$ is feasible, if $x+\alpha d$ is feasible for all $\alpha>0$ that is small enough
- Descent direction $d$ : $d$ is feasible, and $\nabla f(x)^{\top} d<0$


## Algorithm framework

(1) Start at a feasible solution $x^{0}$
(2) Generate a sequence of feasible solutions $x^{k+1}=x^{k}+\alpha_{k} d^{k}$

- $d^{k}$ is a feasible and descent direction
- $\alpha_{k}$ is chosen so that $f\left(x^{k}+\alpha_{k} d^{k}\right)<f\left(x^{k}\right)$
(3) Stepsize rule on $\alpha_{k}$ :
- Armijo criterion
- Constant step size $\alpha_{k}=1$

Q: How to choose an initial feasible solution $x^{0}$ ?
A: When $C$ is a polyhedron, i.e., defined by systems of linear equations/inequalities, we can find one by solving a linear program

## Conditional gradient method (Frank Wolfe)

## Conditional gradient method

A straghtforward way to obtain a descent direction:

$$
\min \nabla f\left(x^{k}\right)^{\top}\left(x-x^{k}\right) \text { s.t. } x \in C
$$

The optimal solution $\bar{x}, d^{k}=\bar{x}-x^{k}$

- $\bar{x}$ will always be on the boundary of $C$
- Makes sense only when this problem is much easier to solve than the original problem. E.g., $f$ is nonlinear, $C$ is a polyhedron
- Convergence could be very slow: sublinear convergence, $\lim _{k \rightarrow \infty} \frac{\left\|x^{k+1}-x^{*}\right\|}{\left\|x^{k}-x^{*}\right\|}=1$ in some cases
- Works well for problems with a low requirement on solution accuracy


## Fundamental theorem in gradient projection

## Theorem

$\Omega$ is a nonempty closed convex set, let $x^{*} \in \Omega$
(a) $x^{*} \in \operatorname{argmin}_{x \in \Omega} f(x) \Rightarrow P_{\Omega}\left(x^{*}-\lambda \nabla f\left(x^{*}\right)\right)=x^{*}, \forall \lambda>0$
(b) If $P_{\Omega}\left(x^{*}-\lambda \nabla f\left(x^{*}\right)\right)=x^{*}$ for some $\lambda>0$, and $f$ is convex, then $f\left(x^{*}\right)=\min _{x \in \Omega} f(x)$

- Note something interesting here: If $P_{\Omega}\left(x^{*}-\lambda \nabla f\left(x^{*}\right)\right)=x^{*}$ for some $\lambda>0$, then $P_{\Omega}\left(x^{*}-\lambda \nabla f\left(x^{*}\right)\right)=x^{*}$ for all $\lambda>0$
- Condition (a) is called gradient projection optimality condition (GPOC)
- GPOC is a generalized definition for stationary point, $\nabla f(x)=0$


## Gradient projection algorithm

(1) GPOC can be seen as a fixed point structure: $F\left(x^{*}\right)=x^{*}$, where $F\left(x^{*}\right)=P_{\Omega}\left(x^{*}-\lambda \nabla f\left(x^{*}\right)\right)$
(2) GPOC can be seen as steepest descent + projection, which is intuitive!
Recall: Stepsize selection problem? $\phi(\alpha)=f\left(x_{k}+\alpha p_{k}\right)$ Wolfe condition

- $\phi(\alpha) \leq \phi(0)+c_{1} \phi^{\prime}(0) \alpha, 0<c_{1}<1$
- $\phi^{\prime}(\alpha) \geq c_{2} \phi^{\prime}(0), 0<c_{2}<1$

Q: What is the problem? $\phi(\lambda)$ here is nonsmooth! So $\phi^{\prime}(\lambda)$ is not available! We cannot use Wolfe condition!

## Armijo backtracking algorithm

Armijo criterion:

$$
f\left(x_{k}\left(\beta^{m} \lambda\right)\right) \leq f\left(x_{k}\right)+c \nabla f\left(x_{k}\right)^{\top}\left(x_{k}\left(\beta^{m} \lambda\right)-x_{k}\right)
$$

Choose the smallest $m$ that the above holds, $\beta \in(0,1)$

- Choose an initial $\lambda$
- Try points $P_{\Omega}\left(x_{k}-\beta^{m} \lambda \nabla f\left(x_{k}\right)\right)$, for $m=0,1, \ldots$
- Stop when sufficient decrease holds


## Theorem

(a) There always exists a qualifying stepsize that satisfies Armijo criterion
(b) Gradient projection algorithm converges to a generalized stationary point

## Rate of convergence

Consider a strictly convex quadratic function $f(x)=\frac{1}{2} x^{\top} Q x-b^{\top} x$, let $x^{*}$ be the unique minimizer of $f$ over $\Omega$. Consider using a constant step size $s$, then:

$$
\begin{aligned}
\left\|x^{k+1}-x^{*}\right\| & =\left\|\left[P_{\Omega}\left(x^{k}-s \nabla f\left(x^{k}\right)\right)\right]-\left[P_{\Omega}\left(x^{*}-s \nabla f\left(x^{*}\right)\right)\right]\right\| \\
& \leq\left\|\left(x^{k}-s \nabla f\left(x^{k}\right)\right)-\left(x^{*}-s \nabla f\left(x^{*}\right)\right)\right\| \\
& =\left\|(I-s Q)\left(x^{k}-x^{*}\right)\right\| \\
& \leq \max \{|1-s m|,|1-s M|\}\left\|x^{k}-x^{*}\right\|
\end{aligned}
$$

where $m$ and $M$ are the smallest and largest eigenvalues of $Q$.

## Concern on gradient projection

(1) Convergence rate same as steepest descent, which is slow!
(2) Gradient projection is still hard! Projection operator is really heavy

- Work well on REALLY simple constraints, e.g., box constraints, where the projection is easy


## Proximal point

## Proximal point

$$
\operatorname{prox}_{P}(x)=\underset{y}{\operatorname{argmin}} \frac{1}{2}\|x-y\|_{2}^{2}+P(y)
$$

where:

- $P(y)$ is an extended value convex function: can take value $+\infty$ and $-\infty$
- "=" is well-defined because of strong convexity


## Examples:

- Indicator function: $\mathbf{1}_{C}(x)=0$ if $x \in C$, and $+\infty$ if $x \notin C$

Q: What is $\operatorname{prox}_{1_{C}}(x) ? \operatorname{proj}_{C}(x)$ !

- $P(x)=\frac{\mu}{2}\|x\|_{2}^{2}$

Q: What is $\operatorname{prox}_{P}(x) ? \frac{1}{1+\mu} x$, shrink $x$ towards origin

## Proximal point: extended gradient projection

## A decomposed unconstrained problem

$$
\min h(x)=f(x)+P(x)
$$

where $f$ is smooth, and $P$ is convex
Extended projection gradient: iterative alternating between proximal point and gradient direction

$$
x_{k+1}=\operatorname{prox}_{\alpha_{k} P} P\left(x_{k}-\alpha_{k} \nabla f\left(x_{k}\right)\right)
$$

## Theorem

If $f$ is convex, $P$ is convex, then $x^{*} \in \operatorname{argmin}_{x} f(x)+P(x)$ if and only if $x^{*}=\operatorname{prox}_{v P}\left(x^{*}-v \nabla f\left(x^{*}\right)\right.$

## Proximal point method

Q: How about solving constrained optimization problem?

$$
\min _{x \in C} f(x) \Leftrightarrow \min f(x)+\mathbf{1}_{C}(x)
$$

Advantage of proximal point method:

- Allow us to solve nonsmooth function minimization at a linear rate
- For more information, check out Convex Optimization by Boyd


## Next time

Penalty functions for general constrained optimization Chapter 17 NW book

