#### Spies and Revolutionaries

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Ex:  $P_9$  is spy-friendly. Consider m = 3, r = 13, s = 4.  $r \rightarrow r r \leftarrow r r \rightarrow r r \leftarrow r r \leftarrow r r \leftarrow r$  $s \quad s \quad s \quad s \quad s \quad s$ 

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